

# Differential Search Algorithm For The Solution of Economic Dispatch Problem

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**Abstract:** This paper deals with the solution of Economic Dispatch problem of power system by using Differential Search Algorithm. The prime objective of this problem is to schedule the load, between the generators so that total cost of generation of electricity is minimum. Differential Search algorithm(DSA) is based upon Brownian –like irregular movement of the life forms to move starting with one place then onto the next. Scientific model of DSA depicts relocation of living creatures, eu-social, sub-social or pre-social during the time to look the current sustenance ranges since it changes intermittently amid the year. To demonstrate the capability of DSA it is tested on 57-unit large scale power framework.

**Keywords:** Economic load dispatch, Differential Search algorithm, Brownian motion.

## I. INTRODUCTION

The Economic Dispatch (ED) issue is a deliberate errand amongst the most fundamental optimization issues in time of the front line power system [1-4]. ED is characterized as the way toward allotting power values to the thermal units in such a way, so that the load supplied altogether will be at minimum cost while fulfilling different constraints. Throughout the years, numerous endeavours have been made to take care of this issue, fusing various types of constraints or different objectives, through different numerical programming and optimization measures. The traditional strategies incorporate Lambda iteration strategy, base point, gradient strategy [1]. These conventional established dispatch algorithms are based on Lagrangian multiplier technique for which the cost curve must be incremental in nature.

In power system there is always losses in the transmission network. For better and practical feasible solutions losses also must take into the ED problem [5-8].

The resulting ED is a non-convex, non-linear optimization problem, which is a challenging task and cannot be solved by the traditional methods. Researcher did several investigations on ELD problem to provide better solutions. As it would result in saving of money at generation end. In the last decade researchers look to the nature to find optimal

solutions of some practical problems of this physical world. In this same effort a novel nature inspired algorithm is used to solve highly constrained problem of power system which is economic dispatch problem. The solution is elegant, considering complexity of the economic load dispatch problem of large scale power system. The convergence rate is fast and also not sensitive to the controlling of parameters.

## II. PROBLEM FORMULATION OF ELD

### A. Minimization of Cost:

The main objective of ELD is to minimize the cost. Which is given by mathematically as per the equation 1.

$$\min F_t = \sum_{i=1}^m F_{th}(P_i) \quad (1)$$

$F_{th}(P_i)$  is thermal cost. The cost associated with thermal power generation can be represented as :

$$F_{th}(P_i) = (a_i P_i^2 + b_i P_i + c_i) (\$/hr.) \quad (2)$$

Where  $a_i$ ,  $b_i$ , and  $c_i$  are the fuel cost coefficients of  $i^{th}$  thermal unit.

Considering valve point loading (VPL) effect thermal power generation cost depicted as:

$$F_{th}(P_i) = a_i P_i^2 + b_i P_i + c_i + |d_i \sin (e_i (P_i^{min} - P_i))| (\$/hr.) \quad (3)$$

Where,  $d_i$  and  $e_i$  are fuel cost coefficients corresponding to VPL effect;  $m$  is the number of thermal units.

### B. Practical Constraints

Power balance constraints:

$$\sum_{i=1}^m P_i = P_D + P_L \quad (4)$$

Where  $P_D$  is the total demand ;  $P_L$  is the transmission network losses, can be represented using B- coefficients:

$$P_L = \sum_{i=1}^m \sum_{j=1}^m P_i B_{ij} P_j + \sum_{i=1}^m B_{oi} P_i + B_{oo} \quad (5)$$

Generation limits constraints:

$$P_{i \min} \leq P_i \leq P_{i \max}$$

$P_{i \min}$ : Minimum power of  $i^{\text{th}}$  generator

$P_{i \max}$ : Maximum power of  $i^{\text{th}}$  generator

### III. DIFFERENTIAL SEARCH ALGORITHM

Differential search algorithm is developed by Pinar Civicioglu [9] for solution of optimization problem. It simulates the random walk like Brownian motion used by an organism to migrate. In the nature several species shows migration behaviour throughout the year. In nature capacity and efficacy of food varies from time to time. Species move from one place to another since reduction in the food capacity of one habitat and move towards where the efficiency of food is more than previous habitat. Many species like species of birds, honeybees and whales have a periodically cycle of immigration. Movement of these species is carried out by super-organism having large number of individuals. Then this super-organism changes its position in search of better food options.

There are many algorithms that are based upon superorganisms like PSO, Cuckoo search, artificial bee colony and Ant colony. Raptorial living beings can control the fertility of the site which they want to migrate. In DS algorithm first the random population is created and then this population change its position in order to find the local minima. During this process artificial-super-organism tests that the selected position is suitable or not temporarily during the migration. If the selected position is suitable to not moving further for a short duration during the migration, then members of species stay there for a while and continue their migration from this position onwards.

Pseudo-code for DS algorithm is shown in Fig-1. In DS algorithm initialize process is defined as

$$x_i = lb + rand * (ub - lb) \quad (6)$$

here  $lb$  is the lower bound of the problem and  $ub$  is the upper bound of the problem.

Artificial-organisms are defined by  $X_i = [x_{ij}]$  and artificial-superorganism made up of artificial-organisms and is defined as superorganism $_g = [X_i]$ .

In DS algorithm mechanism of finding a stopover site may be defined as Brownian-like random walk. Randomly selected individuals of artificial-organisms move towards target of donor  $= [X_{random\_shuffling(i)}]$  for discovery of stopover sites. Random\_shuffling function of given algorithm

changes the order of number of elements in set  $i = \{1, 2, 3, \dots, N\}$ . Size of changes of position of members of artificial-organisms is controlled by scale value that is given in the line 5 of Fig-1. Scale value is given under

$$Scale = randg[2.rand_1].(rand_2 - rand_3) \quad (7)$$

Scale value is produced by gamma random numbers (i.e., randg) controlled by uniform random number generator (i.e., rand) in the range of [0 1] together. Scale value allows to change the direction of artificial-superorganism radically in the habitat.

In DS algorithm stopover site is given in equation 8

$$Stop\ oversite = superorganism\ scale. *(donor - superorganism) \quad (8)$$

where donor = Superorganism Random shuffling

The members of artificial-organisms of the superorganism to participate in the search process of stopover site are determined by a random process. Structure of this random process is given in Fig-1 (line 8-29). If for some reason any element of stopover site goes above the upper limit or below the lower limit then this element is deferred to another position of the habitat (Fig-1, line 31-33).

DS algorithm has only two control parameters  $p_1$  and  $p_2$  (Fig-1, line 7). Most appropriate values of  $p_1$  and  $p_2$  are  $p_1 = 0.3.rand$  and  $p_2 = 0.3.rand$  provides the best solution of ELD problem.

### IV. IMPLEMENTATION OF DS ALGORITHM TO ELD PROBLEM

In this article differential search algorithm is successfully implemented to different ELD problems. There are many steps for ELD implementation which are given below

**Step 1.0-** In step 1 the population of superorganisms which represent the power dispatched by different generators is generated by the equation (9) and it is given under

$$P_i = P_i^{min} + rand * (P_i^{max} - P_i^{min}) \quad (9)$$

$P_i$  = Vector contains population of super-organisms which also represents the different value of power dispatched by generators

$P_{i \min}$  = Minimum value of power that can be generated by the  $i^{\text{th}}$  generator.

$P_{i \max}$  = Maximum value of power that can be generated by the  $i^{\text{th}}$  generator.

rand = any random number between 0 and 1

**Step 2.0-** In step 2 different possible generations from different generators ranked on the basis of their fitness i.e. to find the minimum cost.

**Step 3.0** – In step 3 donor point is generated with the help of superorganism as shown in equation (10) and (11) in order to discover minimum cost.

$$donor(i,:) = superorganism(rs(i,:),) \quad (10)$$

where rs is randomly permuted vector.

**Step 4.0-** In this step a scale value is generated to control the change of size occurred in the position of member of artificial-organisms. The scale value is generated by using gamma random number generator which is given under.

$$scale = rang(2, rand_1). (rand_2 - rand_3) \quad (11)$$

**Step 5.0-** In step-3 the points obtained above are further improve to find the minimum cost by the stopover-site process, in DSA stop-over-site is defined as the where a particular super-organism stays for a while and then choose its habitat according to the fitness of habitat.

$$stopoversite = superorganism\_scale.(donor - superorganism) \quad (12)$$

**Step 6.0-** If any one element of the stopover-site goes beyond the search space then it is randomly placed to another position which is described according to equation 13

$$\text{if } \begin{matrix} stopoversite < P_i^{min} \text{ or} \\ stopoversite > P_i^{max} \end{matrix} \text{ then}$$

$$stopoversite = P_i^{min} + rand * (P_i^{max} - P_i^{min}) \quad (13)$$

Evaluate stopover-site matrix, create the new superorganism and create the new Superorganism vector.

**Step 7.0-** in step 6 if minimum cost is found then this process stops here if minimum value of cost is not found in this step then it repeats steps 1 to 6.

### V. Results and Discussion:

In this section performance of applied algorithm is judged using two test system of ELD problem with 13 and 40 thermal generators. The test cases as mentioned for ED problem is solved using DS algorithm.

DSA potential is verified on a very large scale power system of 300 buses with 57 generating unit. The cost matrix co-

efficient of this test system is given from [10]. The power demand is set 235235.85 MW. The optimum power output achieved by DSA approach is listed in Table-1. Over repeated 20 trial, optimal solution in terms of cost is **566797.446268 \$/hr**. Comparison of convergence result with Augmented Lagrange Hopfield Network (ALHN) [10], particle swarm optimization (PSO) [10], differential evolution (DE) [10] and GSA [11] is presented in Table-2. The convergence graph for 57-unit system is depicted in “Fig.2”.

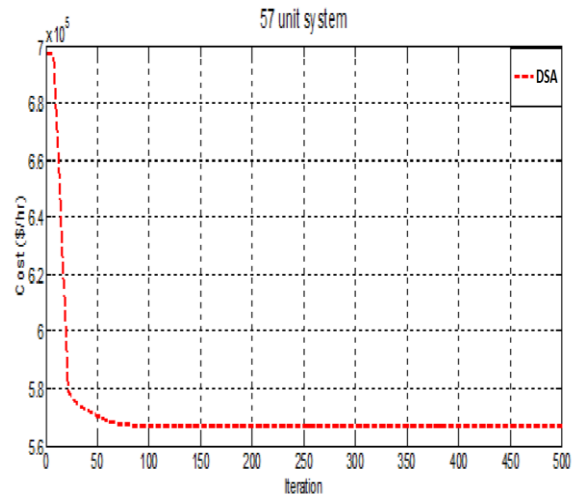


Figure 1 Convergence Characteristics of DSA

Table 1 Comparison of cost obtained by DSA for 57-unit test system

Method	Generation Cost(\$/hr)
DE [10]	587049.0000
PSO [10]	585276.0000
ALHN [10]	584932.0000
GSA [11]	566797.4932
<b>DSA</b>	<b>566797.4462</b>

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Algorithm: Differential search algorithm
N: The size of population where i = {1,2,3,4,.....N}.
D: The dimension of the problem.
G: Number of maximum generation.
1: superorganism = initialize() , where superorganism = [ Artificial-
organismi]
2: yi = Evaluate (Artificial-organismi)
3: for cycle =1:G do
4: donor = Superorganism Random_shuffling(i)
5: Scale = randg[2.rand1].(rand2-rand3)
6: stopoversite=superorganism+scale.(donor-superorganism)
7: p1=0.3.rand4 and p2=0.3.rand5
8: if rand6 < rand7 then
9: if rand8 < p1 then
10: r=rand(N,D)
11: for counter1=1:N do
12: r(counter1,:)=r(counter1,:)<rand9
13: end for
14: else
15: r=ones(N,D)
16: for counter2=1:N do
17: r(counter2,randi(D))=r(counter2,randi(D))<rand10
18: end for
19: end if
20: else
21: r=ones(N,D)
22: for counter3=1:N do
23: d = randi(D,1,[p2.rand.D])
24: for counter4 = 1:size(d) do
25: r(counter3,d(counter4)) =0
26: end for
27: end for
28: end if
29: individualsi,j ← ri,j > 0 | I ∈ I, J ∈ [1,D]
30: stopoversite(individualsi,j) :=superorgansim(individulasi,j)
31: if stopoversitei,j < lowi,j or stopoversitei,j > upi,j
then
32: stopoversitei,j :=lowj+rand.(upj-lowj)
33: endif
34: Ystopoversite,i =evaluate(stopoversitei)
35: Ysuperorganism,i := f(x) = { Ysuperorganism,i if Ystopoversite,i < Ysuperorganism,i
Ysuperorganism,i,i, else
36: ArtificialOrgaismi := f(x) = { Stopoversitei if Ystopoversite,i < Ysuperorganism,i
ArtificialOrganismi, else
37: end for

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Figure 2: Pseudo Code of Differential Search Algorithm

Table 2 Optimal generation of 57-unit system with power demand 23525.85 MW.

Unit (MW)	DSA	Unit (MW)	DSA
P <sub>1</sub>	318.0000	P <sub>30</sub>	255.0000
P <sub>2</sub>	233.0000	P <sub>31</sub>	900.0000
P <sub>3</sub>	435.0000	P <sub>32</sub>	291.591
P <sub>4</sub>	102.0000	P <sub>33</sub>	467.0000
P <sub>5</sub>	99.0000	P <sub>34</sub>	488.0000
P <sub>6</sub>	1640.0000	P <sub>35</sub>	144.0000
P <sub>7</sub>	204.0000	P <sub>36</sub>	126.0000
P <sub>8</sub>	238.0000	P <sub>37</sub>	396.0000
P <sub>9</sub>	591.0000	P <sub>38</sub>	604.561
P <sub>10</sub>	71.0000	P <sub>39</sub>	1028.0000
P <sub>11</sub>	184.0000	P <sub>40</sub>	198.0000
P <sub>12</sub>	87.0000	P <sub>41</sub>	316.0000
P <sub>13</sub>	316.0000	P <sub>42</sub>	495.0000
P <sub>14</sub>	183.0000	P <sub>43</sub>	157.0000
P <sub>15</sub>	174.0000	P <sub>44</sub>	348.0000
P <sub>16</sub>	342.0000	P <sub>45</sub>	425.0000
P <sub>17</sub>	71.0000	P <sub>46</sub>	31.0000
P <sub>18</sub>	300.0000	P <sub>47</sub>	684.0000
P <sub>19</sub>	1800.0000	P <sub>48</sub>	38.0000
P <sub>20</sub>	1020.0000	P <sub>49</sub>	248.0000
P <sub>21</sub>	403.0000	P <sub>50</sub>	347.239
P <sub>22</sub>	1677.0000	P <sub>51</sub>	600.0000
P <sub>23</sub>	360.0000	P <sub>52</sub>	174.0000
P <sub>24</sub>	231.0000	P <sub>53</sub>	1098.0000
P <sub>25</sub>	105.453.0000	P <sub>54</sub>	595.0000
P <sub>26</sub>	382.0000	P <sub>55</sub>	470.0000
P <sub>27</sub>	212.0000	P <sub>56</sub>	42.0000
P <sub>28</sub>	257.0070	P <sub>57</sub>	6.0000
P <sub>29</sub>	518.0000		
<b>Minimum cost (\$/hr)</b>		<b>566797.4462</b>	

VI. Conclusion

DSA has potential to solve the Economic Load Dispatch with considering the practical operating constraints of the modern power system. The results are better than any heuristic algorithm applied for ELD problem. A large power system network of 300 buses with 57 thermal power units is considered as a test case, DSA has solved this large power system with better results. For future perspective we can say that applied algorithm can be applied further in many power system applications like Economic

Emission Dispatch problem, Dynamic Economic Dispatch problem, Optimal Power Flow, Unit commitment.

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